

Hypergraphs with no tight cycles

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UCL

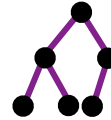
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Introduction

Fact. The max number of edges in a graph on n vertices with no cycles is $n-1$.

(Maximisers are trees.)



An r -graph (= r -uniform hypergraph) is a hypergraph whose edges consist of r vertices.

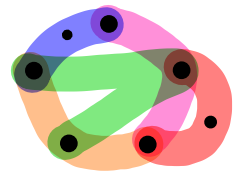


Question. What is the maximum possible number of edges in an r -graph on n vertices which has no cycles?

The answer depends on the definition of a cycle.

We consider three notions: Berge cycles, loose cycles and tight cycles.

Berge cycles

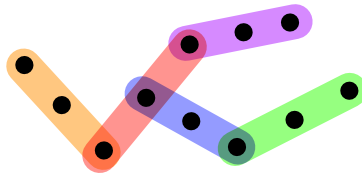


A Berge cycle is a sequence $e_1, v_1, \dots, e_k, v_k$ where e_1, \dots, e_k are distinct edges, v_1, \dots, v_k are distinct vertices, and $v_i \in e_i \cap e_{i+1}$ for $i=1, \dots, k-1$ and $v_k \in e_k \cap e_1$.

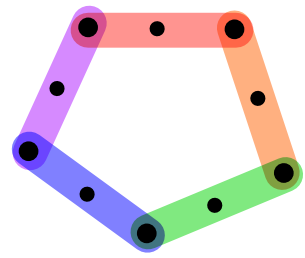
Claim. The maximum number of edges in an r -graph on n vertices with no Berge cycles is $\lfloor \frac{n-1}{r-1} \rfloor$.

Proof of upper bound. Not hard to show: no Berge cycle \Rightarrow edges can be ordered as e_1, \dots, e_m s.t. $|e_i \cap (e_1 \cup \dots \cup e_{i-1})| \leq 1 \quad \forall i$.

Proof of lower bound.



Loose cycles



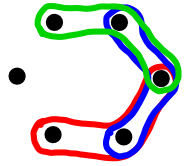
A loose cycle is a sequence e_1, \dots, e_k of edges s.t. consecutive pairs of edges (and e_k and e_1) have exactly one vertex in common, and non-consecutive pairs are disjoint.

Theorem (Frankl-Füredi '87). The maximum number of edges in an r -graph on n vertices with no loose cycles is $\binom{n-1}{r-1}$ (for n large).

Actually, FF show that $> \binom{n-1}{r-1}$ edges $\Rightarrow \exists$ loose triangle.
(cycle of length 3)

Proof of lower bound. Consider the r -graph with vertices $[n] := \{1, \dots, n\}$ whose edges are all sets of r vertices containing 1.

Tight cycles.



A tight cycle is a sequence v_1, \dots, v_k of distinct vertices s.t. $\{v_i, \dots, v_{i+r-1}\}$ is an edge $\forall i$ (indices are modulo k).

$f_r(n) = \max$ number of edges in an r -graph on n vertices with no tight cycles.

Lower bounds

Observation. $f_r(n) \geq \binom{n-1}{r-1}$.

Proof. Consider r -graph \mathcal{H} with $V(\mathcal{H}) = [n]$ and $E(\mathcal{H}) = \{r\text{-sets containing } 1\}$.

Conjecture (Sós, Verstraëte '10s). $f_r(n) = \binom{n-1}{r-1}$ for large n .

Huang-Ma '19. The conjecture is false (for $r \geq 3$).

There exists $c \in (1, 2)$ s.t. $f_r(n) \geq c \binom{n-1}{r-1}$.

B. Janzer '20. Let $r \geq 3$. There exists $c > 0$ s.t. $f_r(n) \geq \frac{cn^{r-1} \log n}{\log \log n}$.

Upper bounds

A result of Erdős ('64) implies $f_r(n) = O(n^{r-2^{r-1}})$.

An unpublished result of Verstraëte implies $f_3(n) = O(n^{3/2})$.

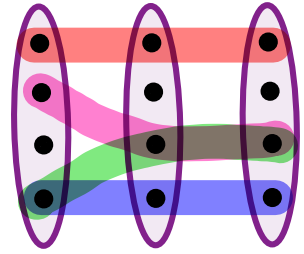
Sudakov-Tomon ('21). Let $r \geq 3$. Then $f_r(n) \leq n^{r-1} e^{O(\sqrt{\log n})}$.

In particular, $f_r(n) = n^{r-1+o(1)}$.

Theorem (L. '21+). Let $r \geq 3$. Then $f_r(n) \leq O(n^{r-1} (\log n)^5)$.

(This is tight up to a factor of $O((\log n)^4 \log \log n)$.)

r -partite r -graphs



Claim. Let \mathcal{H} be an r -graph on m edges. There is an r -partite subhypergraph $\mathcal{H}' \subseteq \mathcal{H}$ on $\geq \frac{r!}{r^r} \cdot m$ edges.

Proof. Let $\{A_1, \dots, A_r\}$ be a partition of $V(\mathcal{H})$, chosen uniformly at random, and let \mathcal{H}' be the subhypergraph of \mathcal{H} obtained by keeping edges e with $|e \cap A_i| = 1 \forall i$.

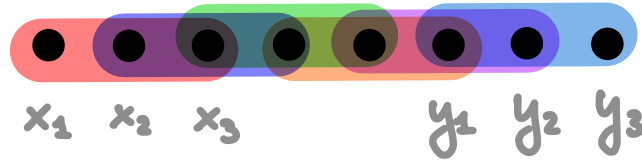
For $e \in \mathcal{H}$, $\mathbb{P}(e \in \mathcal{H}') = \frac{r!}{r^r} \Rightarrow \mathbb{E}(e(\mathcal{H}')) = \frac{r!}{r^r} \cdot m \Rightarrow$ suitable \mathcal{H}' exists.

Aim. Let \mathcal{H} be an r -partite r -graph, with partition $\{A_1, \dots, A_r\}$, on n vertices with $\geq dn^{r-1}$ edges, where $d \geq c \cdot (\log n)^5$. Then \mathcal{H} contains a tight cycle.

Notation

We represent edges $e \in E$ as (x_1, \dots, x_r) where $x_i \in A_i$.

For edges $e = (x_1, \dots, x_r)$ and $f = (y_1, \dots, y_r)$, a tight path from e to f is a tight path (z_1, \dots, z_ℓ) where $z_i = x_i$ for $i \in [r]$ and $z_{\ell-r+i} = y_i$ for $i \in [r]$.



Expanders

r -partite r -graph, $\geq dn^{r-1}$ edges

Lemma (Sudakov-Tomon '21). There is a subhypergraph $\mathcal{G} \subseteq \mathcal{H}$ s.t. for every edge $e \in \mathcal{G}$ and small $B \subseteq V(\mathcal{G})$ (which is disjoint of e): for almost every edge $f \in \mathcal{G}$ there is a short tight path from e to f that avoids B .

Proof ideas.

- * Define r -line graphs, which are graphs associated with r -partite r -graphs.
- * Define notions of density and expansion for r -line-graphs.

Call such \mathcal{G} an expander.

Connecting edges in expanders.

Theorem (Sudakov-Tomon '21). Let \mathcal{G} be an expander. Then

- * Either for every disjoint edges e, f there is a tight path from e to f ,
- * Or there is a small dense subhypergraph.

They proceed by a density increment argument.

Theorem (L. 21+). Let \mathcal{G} be an expander. Then \mathcal{G} contains a tight cycle.

In fact, we show that there is a tight cycle through almost every two edges.

Robust reachability

Let \mathcal{G} be an expander with m edges. So, for every $e \in E(\mathcal{G})$ and $B \subseteq V(\mathcal{G})$ with $|B| \leq \frac{\ell^2}{\varepsilon}$, which is disjoint of e , for at least $(1-\varepsilon)m$ edges $f \in E(\mathcal{G})$: there is a tight path from e to f of length $\leq \ell$ that avoids B .

Lemma (L.). In \mathcal{G} as above, for every edge e there is a set $F(e) \subseteq E(\mathcal{G})$ with $|F(e)| \geq (1-\varepsilon)m$, and paths $\mathcal{P}(e,f)$, for $f \in F(e)$, of length $\leq \ell$ from e to f , s.t. no vertex (except for those in e) appears in $> \frac{\varepsilon m}{\ell}$ paths $\mathcal{P}(e,f)$.

Proof of lemma

* Let $F(e) \subseteq E(\mathcal{G})$ be maximal s.t. \exists paths $\mathcal{P}(e,f)$ of length $\leq l$ from e to f , for $f \in F(e)$, s.t. no vertex appears in $> \frac{\epsilon m}{l}$ paths $\mathcal{P}(e,f)$.

* Let B be the set of vertices in exactly $\frac{\epsilon m}{l}$ paths $\mathcal{P}(e,f)$.

$$\text{Then } |B| \cdot \frac{\epsilon m}{l} \leq l \cdot |F(e)| \leq lm \Rightarrow |B| \leq \frac{l^2}{\epsilon}.$$

* By expansion, there is a subset $F'(e) \subseteq E(\mathcal{G})$ of size $\geq (1-\epsilon)m$, and paths $\mathcal{P}(e,f)$ for $f \in F'(e)$, of length $\leq l$ from e to f , that avoids B .

* If $|F(e)| < (1-\epsilon)m$, let $f' \in F'(e) \setminus F(e)$.

The set $F(e) \cup \{f'\}$ with paths $\mathcal{P}(e,f)$ for $f \in F(e)$ and $\mathcal{P}(e,f')$ for f' , contradicts maximality of $F(e)$.

Finding tight cycles

Let \mathcal{G} be an expander with m edges. So, for every $e \in E(\mathcal{G})$ and $B \subseteq V(\mathcal{G})$ with $|B| \leq \frac{\ell^2}{\varepsilon}$, which is disjoint of e , for at least $(1-\varepsilon)m$ edges $f \in E(\mathcal{G})$: there is a tight path from e to f of length $\leq \ell$ that avoids B .

Theorem (L. 21+). Let \mathcal{G} be as above. Then \mathcal{G} contains a tight cycle.

Lemma (L.). In \mathcal{G} as above, for every edge e there is a set $F(e) \subseteq E(\mathcal{G})$ with $|F(e)| \geq (1-\varepsilon)m$, and path $\mathcal{P}(e,f)$ of length $\leq \ell$ from e to f , for $f \in F(e)$, s.t. no vertex (except for those in e) appears in $> \frac{\varepsilon m}{\ell}$ paths $\mathcal{P}(e,f)$.

Proof of theorem

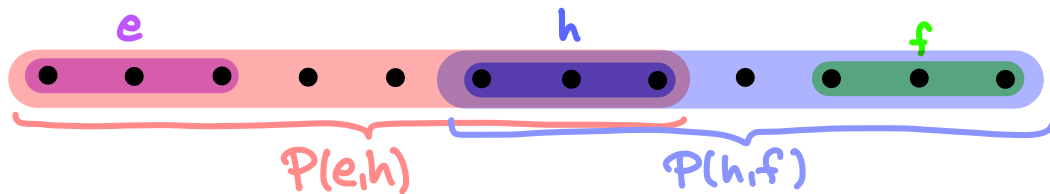
- * For every $e \in E(\zeta)$, let $F(e) \subseteq E(\zeta)$ be a set of size $\geq (1-\varepsilon)m$, and let $\mathcal{P}(e,f)$ be a tight path of length $\leq \ell$ from e to f , for $f \in F(e)$, s.t. no vertex appears in $> \frac{\varepsilon m}{\ell}$ paths $\mathcal{P}(e,f)$ with $f \in F(e)$.
- * For $f \in E(\zeta)$, let $B(f)$ be the set of vertices appearing in $\geq \frac{\varepsilon m}{\ell}$ paths $\mathcal{P}(e,f)$.
Then $|B(f)| \cdot \frac{\varepsilon m}{\ell} \leq \ell m \Rightarrow |B(f)| \leq \frac{\ell^2}{\varepsilon}$.
- * By expansion, for $\geq (1-\varepsilon)m$ edges e there is a path $Q(f,e)$ of length $\leq \ell$ from f to e that avoids $B(f)$.

Many good pairs

Claim. There are at least $(1-7\epsilon)m^2$ pairs (e,f) with $e,f \in E(\mathcal{G})$ s.t.:

(a) $Q(f,e)$ is defined,

(b) there are at least $\frac{m}{2}$ edges h s.t. $P(e,h)$ and $P(h,f)$ are defined and the concatenation $P(e,h)P(h,f)$ is a tight path from e to f .



Proof of claim

Let $\mathcal{D}(e)$ be the digraph with vertices $E(\mathcal{G})$ where hf is an edge whenever $P(e,h), P(h,f)$ are defined and $P(e,h)P(h,f)$ is a tight path.

For every h s.t. $P(e,h)$ is defined, h has out-degree at least:

$$\underbrace{\#\{f \in E(\mathcal{G}) : P(h,f) \text{ is defined}\}}_{\geq (1-\varepsilon)m} - \sum_{\substack{v \in V(P(e,h)) \\ v \neq h}} \underbrace{\#\{f : v \text{ is in } P(h,f)\}}_{\leq \frac{\varepsilon m}{\ell}} \geq (1-2\varepsilon)m.$$

$\Rightarrow \mathcal{D}(e)$ has $\geq (1-\varepsilon)m \cdot (1-2\varepsilon)m \geq (1-3\varepsilon)m^2$ edges.

$\Rightarrow \mathcal{D}(e)$ has $\geq (1-6\varepsilon)m$ vertices f with in-degree $\geq \frac{m}{2}$.

\Rightarrow (b) is satisfied for $\geq (1-6\varepsilon)m^2$ pairs (e,f) .

End of proof of theorem

Let $e, f \in E(\mathcal{G})$ satisfy:

(a) $Q(f, e)$ is defined,

(b) there are at least $\frac{m}{2}$ edges h s.t. $P(e, h), P(h, f)$ are defined and the concatenation $P(e, h)P(h, f)$ is a tight path from e to f .

* There are at most $\frac{\epsilon_m}{\ell} \cdot \ell = \epsilon_m$ edges h as in (b) s.t. $P(e, h)$ and $Q(f, e)$ have a vertex in common (which is not in e).

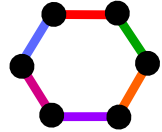
* There are at most $\frac{\epsilon_m}{\ell} \cdot \ell = \epsilon_m$ edges h as in (b) s.t. $P(h, f)$ and $Q(f, e)$ have a vertex in common (which is not in f).

\Rightarrow There are $\geq \frac{m}{2} - 2\epsilon_m \geq \frac{m}{4}$ edges h as in (b) s.t. $P(e, h)P(h, f)Q(f, e)$ is a tight cycle.

Other applications

The same arguments yield the following results (already proved by O. Janzer '20).

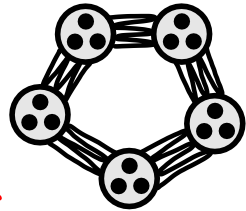
- * Every properly edge-coloured graph on n vertices with $\geq cn(\log n)^5$ edges contains a rainbow cycle.
→ incident edges have distinct colours
↘ all edges have distinct colours



(using tools from Jiang-Methuku-Yepremyan '21; observed by JLMY '21+).

- * Every graph on n vertices with $\geq cn^{2-1/r}(\log n)^{5/r}$ edges contains an r -blowup of a cycle.

↖ vertices replaced by independent sets of size r ; edges replaced by $K_{r,r}$



(using a result by Morris-Saxton '16; connection observed by JLMY).

Open problems

* What is $f_r(n) = \max$ number of edges in an r -graph on n vertices with no tight cycles? Best known bounds: $c_1 n^{r-1} \frac{\log n}{\log \log n} \leq f_r(n) \leq c_2 n^{r-1} (\log n)^5$.
B. Janzer '20 L. 21+

* Question (Conlon '11). Is there $c=c(r)$ s.t. if $\ell > r$ and $r \mid \ell$ then r -graphs on n vertices with no tight cycles of length ℓ have $O(n^{r-1+c/\ell})$ edges?
(Every r -partite r -graph has no tight cycles of length ℓ with $r \nmid \ell$.)

* What is $g(n) = \max$ number of edges in a properly edge-coloured graph on n vertices with no rainbow cycles?

Best known bounds: $c_1 n \log n \leq g(n) \leq c_2 n (\log n)^4$.
Keevash-Mubayi-Sudakov-Verstraëte '07 O. Janzer '20

Thank you!